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# Computation model and higher-dimensional tilings

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## Résumé

Most computability results on subshifts are proved on  $\mathbb{Z}^2$  by embedding arbitrary computations of Turing machines, which classically consists in drawing the space-time diagram of the Turing machine (one dimension being used for space, one dimension being used for time) ; and then, said results generalize to higher dimensions by reducing to the two-dimensional case.

These reductions between the 2D and the  $d$ -dimensional case have left a gap in my understanding of how we compute with subshifts: how do you embed arbitrary computations in a  $d$ -dimensional tiling? For example, if you draw the space-time diagram of a Turing machine on  $\mathbb{Z}^3$ , do you use one dimension for space, and two for time? This appears to suggest a tradeoff between time and space when dealing with  $d$ -dimensional computations.

In this talk, we show that there is, in fact, no tradeoff. By considering "mesh-connected multicomputers" as a computation model, one can implement the sorting of an array of size  $(0,n)^{d-1}$  into a tiling of the  $(0,n)^d$  cube. From this, we can embed  $O(n^{d-1})$  steps of arbitrary word-RAM computations into a tiling of the  $(0,n)^d$  cube.

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